Max-Min Fairness for Beamspace MIMO-NOMA: From Single-Beam to Multi-Beam

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Abstract—With the help of non-orthogonal multiple access (NOMA), the number of connections of the beamspace multiple-input multiple-output (MIMO) systems can be improved with enhanced sum-rate performance, which constitutes beamspace MIMO-NOMA. Thus, most relevant papers focus on improving the system sum rate, which may conflict with the rate loss to weak users. To ensure the achievable rates of weak users, we maximize and analyze the minimal rate of the system in the single-beam case as well as the multi-beam case, where two completely different phenomena are revealed. Particularly, in the single-beam case, the maximized minimal rate of the beamspace MIMO-NOMA always grows rapidly with the signal-to-noise-ratio (SNR), and is larger than that of the beamspace MIMO using orthogonal multiple access (beamspace MIMO-OMA). However, in the multi-beam case, the maximized minimal rate of the beamspace MIMO-NOMA grows slower and slower in the high-SNR region, where it is smaller than that of the beamspace MIMO-OMA. To explain this difference, it is disclosed that the intra-beam interference in the single-beam case is of successive pattern, which is proved to have no limit on the max-min rate. In contrast, the inter-beam interference in the multi-beam case is of mutual pattern, which is proved to restrict the max-min rate to a derived upper bound.

Index Terms—Beamspace MIMO-NOMA, max-min rate, interference, successive pattern, mutual pattern.

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NOMA in a hybrid precoding-based mmWave communication system, and maximized the sum rate of two single-antenna users via joint optimization of power allocation parameters and beamforming vectors. Furthermore, [21] considered a downlink mmWave communication scenario, where the BS and the users are equipped with multiple antennas and single RF chain. NOMA is integrated in this system to serve multiple users, where power allocation parameters, transmit beamforming, and receive beamforming are jointly optimized for sum-rate maximization. Besides, [22] investigated a NOMA-based mmWave communication system with hybrid precoding in the single-beam case, where the joint optimization of power allocation parameters and beamforming vectors is carried out for minimal rate maximization.

In the multi-beam case with the BS equipped with multiple RF chains, [23] firstly explored the beamspace MIMO-NOMA scheme, where NOMA is utilized in the beamspace MIMO based mmWave communication system to support multiple users in each beam. Moreover, power allocation is optimized to maximize its sum rate by the weighted minimum mean square error (WMMSE) algorithm. Similarly, [24] also aimed at improving the sum rate in a multi-beam mmWave communication system integrated with NOMA. To achieve this goal, antenna allocation, user grouping, and power allocation are jointly optimized. Additionally, [25] maximized the sum rate for a NOMA based mmWave communication system, where user clustering is designed by a Stackelberg game-based approach, and the optimal power allocation parameters are acquired in closed-form expressions.

It is worth pointing out that most of the above-mentioned works mainly consider to improve the sum rate, which is an important performance metric extensively explored in NOMA based mmWave communication systems. However, only focusing on the sum rate may incur substantial rate loss to weak users. This is because the system tends to allocate most of the communication resources to the strong users when the sum rate is maximized. In some extreme cases such as that considered in [26], the weak users even cannot be served. Therefore, in order to guarantee the achievable rates for weak users, the important performance metric of user fairness needs to be considered. Correspondingly, we maximize the minimal rate of the system from the perspective of max-min fairness. Note that some researchers tend to set a target data rate for all users to avoid the extreme case in [26]. Under this circumstance, studying the maximized minimal rate still makes sense, since it provides an upper bound for the target data rate.

B. Our Contributions

To ensure the rate performance for all users, we maximize and analyze the minimal rate of beamspace MIMO-NOMA considering the single-beam case as well as the multi-beam case. The corresponding theoretical analysis is challenging, because the interference patterns of the considered system are complicated, especially for the multi-beam case which suffers from both the intra-beam and inter-beam interferences. As a result, the minimal rate maximization problem is hard to solve, and it is difficult to analyze the effects of both the interferences on the maximized minimal rate (max-min rate). To this end, we utilize a bisection-based method for minimal rate maximization, disclose the different patterns of the intra-beam as well as the inter-beam interferences, and comprehensively analyze their different effects on the max-min rate. To the best of our knowledge, we are the first to carry out the max-min fairness analysis for this system in both of the two cases.\footnote{Simulation codes are provided to reproduce the results presented in this paper: http://oa.ee.tsinghua.edu.cn/dailinglong/publications/publications.html.}

1) Different from [22] and [27] which maximize the minimal rate of NOMA-based single-antenna/single-beam systems, we maximize the minimal rate of beamspace MIMO-NOMA in both the single-beam and multi-beam cases via power allocation. Moreover, we point out the difference between the two cases in terms of the inter-beam interference, and extend the single-beam property that the max-min rate is achieved when rates of all NOMA users are equal to one another to the multi-beam case.

2) We analyze the max-min rate of beamspace MIMO-NOMA by comparing it with that of the beamspace MIMO using orthogonal multiple access (beamspace MIMO-OMA) in both cases, and reveal that the comparison results in the two cases are completely different. To be more specific, in the single-beam case, beamspace MIMO-NOMA is always superior to beamspace MIMO-OMA in terms of the max-min rate, and the max-min rate of both systems grow rapidly with the SNR. However, in the multi-beam case, beamspace MIMO-NOMA is inferior to beamspace MIMO-OMA in the high-SNR region, where the max-min rate of the former system grows slower and slower with the increase of SNR.

3) To account for the difference in the two cases, we point out that the interfering patterns of the intra-beam and inter-beam interferences are completely different, and analyze their different effects on the max-min rate. Specifically, the intra-beam interference in the single-beam case is of successive pattern, and we prove that it will not limit the max-min rate. However, the inter-beam interference in the multi-beam case is of mutual pattern, which is proved to largely constrain the max-min rate, while the upper bound is also derived. Note that the above analysis does not rely on specific system design schemes, and can serve as a general framework for the max-min fairness analysis of beamspace MIMO-NOMA.

C. Organization and Notation

Organization: The remainder of the paper is organized as follows. Section II detailedly introduces the system model of beamspace MIMO-NOMA. Section III maximizes the minimal rate of all NOMA users in the single-beam case, which is shown to be always larger than that of the traditional beamspace MIMO-OMA. Moreover, Section IV maximizes the minimal rate of all NOMA users in the multi-beam case, and reveals that it is smaller than the maximized minimal
rate of beamspace MIMO-OMA in the high-SNR region. To explain this difference, Section V points out the different interfering patterns of the intra-beam and inter-beam interferences, and analyze their different effects on the max-min rate. Finally, Section VI draws the conclusion.

Notation: Vectors and matrices are denoted by lower-case and upper-case boldface letters, respectively. Conjugate transpose, Moore-Penrose inversion, and inversion of the matrix and upper-case boldface letters, respectively. Conjugate transpose.

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The single-cell downlink beamspace MIMO-NOMA system is elaborated in this section. To provide a clearer description, the intrinsic sparsity of mmWave channel, the spatial channel can be mathematically described as an.

![Fig. 1. System model of beamspace MIMO-NOMA.](image)

**Fig. 1.** System model of beamspace MIMO-NOMA.
where \( \hat{h}_{m,n}^H \) denotes the corresponding beamspace channel vector of size \( N_{RF} \times 1 \) after beam selection, and \( v_{m,n} \sim \mathcal{CN}(0, \sigma^2) \) denotes the thermal noise. Note that the digital precoding vectors \( \{w_m\} \) are the same for the users served by the same beam. They are calculated by the SVD-based zero-forcing method proposed in [23], which exploits the beamspace channel vectors of all the NOMA users in each beam. With the digital precoding vectors, the NOMA users in the same beam are sorted in the descending order of their beamspace channel vectors of all the NOMA users in each beam.

To decode multi-user signal, SIC is adopted by the users within each beam. Particularly, while carrying out SIC, the \( i \)-th user in the \( m \)-th beam \((i < n)\) will firstly decode the signals for the users with worse effective channel qualities, i.e., the users in the same beam with indices larger than \( i \), including the \( n \)-th user. Then, the \( i \)-th user will successively remove them from its received signal, in the order of \( |S_m| \) - 1, \( \cdots \), \( i+1 \), until its own signal can be detected. Therefore, when the \( i \)-th user is about to decode the \( n \)-th user’s signal, the remaining signal \( \hat{y}_{m,i(n)} \) is given by

\[
\hat{y}_{m,i(n)} = \hat{h}_{m,i}^H \sqrt{p_{m,n}} w_m s_{m,n} + \hat{r}_{m,i}^H \sum_{n'=1}^{|S_m|} \sqrt{p_{m,n'}} w_m s_{m,n'} + v_{m,i},
\]

In equation (6), the first term denotes the \( n \)-th user’s signal to be decoded by the \( i \)-th user, the second term represents the intra-beam interference from users in the \( m \)-th beam with better effective channel qualities (their indices are smaller than \( n \)), and the third term denotes the inter-beam interference from users served by other beams. Based on (6), the signal-to-interference-plus-noise-ratio (SINR) for the \( i \)-th user to decode the \( n \)-th user’s signal in the \( m \)-th beam is given by

\[
\gamma_{n,i}^m = \frac{\| \hat{h}_{m,i}^H w_m \|_2^2 p_{m,n}}{\xi_{n,i}^m + \sigma^2},
\]

where \( \xi_{n,i}^m \) denotes the power of the intra- and inter-beam interferences:

\[
\xi_{n,i}^m = \sum_{n'=1}^{n-1} \| \hat{h}_{m,i}^H w_m \|_2^2 p_{m,n'} + \sum_{n' \neq m}^{\{S_m\}} \| \hat{h}_{m,i}^H w_{m'} \|_2^2 p_{m',n'}.
\]

Thus, the corresponding achievable rate can be expressed as \( R_{n,i}^m = \log_2 (1 + \gamma_{n,i}^m) \).

Moreover, in the \( m \)-th beam, SIC requires successful detection of the \( n \)-th user’s signal by the \( i \)-th user in the same beam for all \( i < n \) \([27], [29], [30]\). Therefore, the corresponding achievable rate \( R_n^m \) is the minimum among a series of rates:

\[
R_n^m = \min_{i=1,2,\cdots,n} R_{n,i}^m = \min_{i=1,2,\cdots,n} \log_2 (1 + \gamma_{n,i}^m).
\]

The minimal rate of all NOMA users is given by:

\[
R_{\min} = \min_{m,n} \{ R_n^m \}.
\]

Based on the equations (7) and (8), it is worth pointing out that the nominator of the SINR is the desired signal power, while the denominator includes the power for the intra-beam and inter-beam interferences, both of which are mainly affected by the power allocation parameters \( \{p_k\} \). This indicates that the rates of the NOMA users are mainly determined by power allocation. As a result, power allocation problems will be formulated to maximize the minimal rate in both the single-beam case and the multi-beam case. Moreover, by formulating the power allocation problems, we are improving the nominator and reducing the denominator of the corresponding SINRs. This is to say, power allocation can increase the desired signal power and mitigate the inter-beam and intra-beam interferences. Accordingly, two bisection-based methods are proposed in our paper for power allocation, and they are able to suppress the inter-beam and intra-beam interferences, which are presented in Sections III and IV, respectively.

III. MAX-MIN FAIRNESS ANALYSIS FOR SINGLE-BEAM BEAMSPACE MIMO-NOMA

The maximized minimal rate of beamspace MIMO-NOMA in the single-beam case is analyzed in this section, where only one RF chain is deployed at the BS. Particularly, the minimal rate of all NOMA users is maximized by formulating and solving a single-beam power allocation problem, which is shown to be always larger than that of the traditional beamspace MIMO-OMA.

A. Problem Formulation

For the single-beam beamspace MIMO-NOMA system, the BS is equipped with only one RF chain, which indicates that only one beam can be generated at a time. Therefore, for the ease of notation, we will omit the index for the beam in this section. Moreover, the system is only influenced by intra-beam interference, and the beamspace channel \( \hat{h}_n \) for the \( n \)-th user after beam selection is a scalar. Simplified from (7), the \( i \)-th user can decode the \( n \)-th user’s signal with an SINR \( \gamma_{n,i} \) given by:

\[
\gamma_{n,i} = \frac{|\hat{h}_n|^2 p_n}{K \sum_{k=1}^{n-1} |\hat{h}_k|^2 p_k + \sigma^2},
\]

where \( p_k \) is the power allocated to the \( k \)-th user.

Given that in the single-beam case, the order of effective channel quality reduces to \( |\hat{h}_1|^2 > |\hat{h}_2|^2 > \cdots > |\hat{h}_K|^2 \). Therefore, for any \( i < j < n \), the SINR for the \( i \)-th user to decode the signal for the \( n \)-th user is larger than
that of the $j$-th user to decode the same signal, which is shown by:

$$\gamma_{n,i} - \gamma_{n,j} = \frac{p_{n}\sigma^2(|\hat{h}_i|^2 - |\hat{h}_j|^2)}{\left(\sum_{k=1}^{n-1} |\hat{h}_k|^2 p_k + \sigma^2\right)\left(\sum_{k=1}^{n-1} |\hat{h}_j|^2 p_k + \sigma^2\right)} > 0.$$  \hspace{1cm} (12)

Note that $\gamma_{n,i} - \gamma_{n,j} > 0$ is due to the channel ordering $|\hat{h}_i|^2 - |\hat{h}_j|^2 > 0$, which is independent of the specific value of the power coefficients $\{p_k\}$. Thus, we can conclude that equation (12) always holds, which has nothing to do with the specific value of power coefficients. Based on equation (12), for the $n$-th user, we have $\gamma_{n,1} > \gamma_{n,2} > \cdots > \gamma_{n,n}$, and we know before power allocation that its achievable rate $R_n$ is actually determined by the rate of decoding its signal at itself:

$$R_n = \min_{i=1,2,\ldots,n} R_{n,i},$$

$$= \min_{i=1,2,\ldots,n} \log_2 (1 + \gamma_{n,i}),$$

$$= \log_2 (1 + \gamma_{n,n}),$$

$$= R_{n,n}. \quad \hspace{1cm} (13)$$

Thus, to maximize the minimal rate, the objective function of the single-beam power allocation problem is formulated by

$$\max_{\{p_n\}} \min_{\{p_n\}} \{R_{n,n}\}, \quad \hspace{1cm} (14)$$

and we only need to consider the following two constraints:

$$p_n \geq 0, \quad \forall 1 \leq n \leq K,$$

$$\sum_{n=1}^{K} p_n \leq P_{\max}, \quad \hspace{1cm} (15)$$

where (15) ensures that each user is assigned non-negative power, and (16) sets the maximum transmit power $P_{\max}$. Then, the corresponding single-beam power allocation problem $P_1$ is given by

$$P_1 : \max_{\{p_n\}} \min_{\{p_n\}} \{R_{n,n}\},$$

s.t. $C_1 : p_n \geq 0, \quad \forall 1 \leq n \leq K,$

$$C_2 : \sum_{n=1}^{K} p_n \leq P_{\max}, \quad \hspace{1cm} (17)$$

where the max-min rate is denoted by $r^*$.  

The above problem $P_1$ is challenging, because the achievable rate $R_{n,n}$ is a non-convex function due to the intra-beam interference, as shown by (11) and (13). To cope with this challenge, we will introduce a bisection-based method in the next subsection to solve this problem.

### B. Minimal Rate Maximization

In this subsection, the max-min rate $r^*$ is acquired by utilizing a bisection procedure. To deal with the non-convex objective function, an auxiliary variable $t$ is introduced to equivalently convert the original single-beam power allocation problem $P_1$ into a new problem $P_1'$:

$$P_1' : \max_t,$$

s.t. $C_1 : p_n \geq 0, \quad \forall 1 \leq n \leq K,$

$$C_2 : \sum_{n=1}^{K} p_n \leq P_{\max},$$

$$C_3 : R_{n,n} \geq t, \quad \forall 1 \leq n \leq K. \quad \hspace{1cm} (18)$$

The above converted problem $P_1'$ resembles the power allocation problem in [27], and the bisection-based method proposed in [27] can also be utilized to solve $P_1'$. Note that solving this single-beam power allocation problem is not our contribution. Our contributions mainly lie in the revealing of the different interfering patterns of the intra-beam and inter-beam interferences, as well as the analysis of their different effects on the max-min rate in the single-beam and multi-beam cases. These will be detailedly presented in the following sections.

To be more specific, we can set the value of $t$ to a certain $t_0$ in a bisection manner, and examine whether $t_0$ can be achieved by the max-min rate $r^*$, which is detailedly described in Algorithm 1. For the examination of $t_0$, the achievable rates for all NOMA users are firstly constrained to be larger than $t_0$, and the total power consumption $\sum_{n=1}^{K} p_n$ is minimized to see whether it is smaller than the maximum transmit power $P_{\max}$. If the minimized power consumption is smaller than $P_{\max}$, then $t_0$ can be achieved, otherwise it cannot be achieved. This brings the following optimization problem $P_1'$:

$$\tilde{P}_1 : \min_{\{p_n\}} \sum_{n=1}^{K} p_n,$$

s.t. $C_1 : p_n \geq 0, \quad \forall 1 \leq n \leq K,$

$$C_2 : R_{n,n} \geq t_0, \quad \forall 1 \leq n \leq K. \quad \hspace{1cm} (19)$$

According to Proposition 2 in [27], the minimized power consumption of the problem $P_1$ is achieved by satisfying all

| Algorithm 1 Bisection-Based Single-Beam Power Allocation Algorithm |
|------------------|------------------|------------------|
| **Input:** beamspace channels $\{\hat{h}_i\}$, noise power $\sigma^2$, total power $P_{\max}$, desirable accuracy $\epsilon$. |
| **Output:** Optimal power allocation parameters $\{p_n^*\}$, max-min rate $r^*$. |
| **1:** Set lower bound $t_L = 0$. |
| **2:** Set upper bound $t_H = \log_2(1 + P_{\max}/\max_{i} |\hat{h}_i|^2)$, where $h_{\max}$ is the maximal value among all $|\hat{h}_i|^2$. |
| **3:** while $t_H - t_L > \epsilon$ do |
| **4:** Set $t_0 = (t_H + t_L)/2$, solve problem $\tilde{P}_1$ via equation (20) to obtain $\{p_n\}$. |
| **5:** if $\sum_{n=1}^{K} p_n \leq P_{\max}$ then |
| **6:** Set $t_L = t_0$, $r^* = t_0$, $\{p_n^*\} = \{p_n\}$. |
| **7:** else |
| **8:** Set $t_H = t_0$. |
| **9:** end if |
| **10:** end while |
the total complexity for solving \(K\) is obtained by solving a series of problems \(\hat{P}_1\), it will be achieved when the rates for all NOMA users are equal to one another [27]. Based on this important finding, the optimal solution of \(\hat{P}_1\) can be obtained by solving the following set of equations:
\[
R_{n,n} = t_0
\]
\[
\Leftrightarrow \log_2 \left( 1 + \frac{\vert \hat{h}_n \vert^2 p_n}{\sum_{k=1}^{n-1} \vert \hat{h}_n \vert^2 p_k + \sigma^2} \right) = t_0
\]
\[
\Leftrightarrow p_n = (2^{t_0} - 1) \left( \sum_{k=1}^{n-1} p_k + \frac{\sigma^2}{\vert \hat{h}_n \vert^2} \right), \tag{20}
\]
and we have \(p_1 = (2^{t_0} - 1) \frac{\sigma^2}{\vert \hat{h}_n \vert^2}\). Based on equations (20), we can obtain the power allocation parameters \(p_2, p_3, \ldots, p_K\) successively from \(p_1\), each with linear complexity \(O(K^3)\) [27]. Since there are altogether \(K\) power allocation parameters, the total complexity for solving \(\hat{P}_1\) is \(O(K^3)\), and the optimal power consumption for problem \(\hat{P}_1\) can be calculated as \(P^* = \sum_{n=1}^{K} n\). By comparing \(P^*\) with \(P_{\text{max}}\), we can examine whether \(t_0\) can be achieved as mentioned before, and the bisection-based method is shown in Algorithm 1, where the lower bound is set to 0, and the upper bound is set to the rate of utilizing full power to serve the strongest user.

According to [31], the bisection-based method will converge to the optimal solution of the original problem \(P_1\) within a desirable accuracy \(\epsilon\) with linear convergence, which also demonstrates the effectiveness of the method, and the number of required iterations is approximately \(O(\log(\frac{1}{\epsilon}))\) [31]. Since in each iteration the computation complexity to solve \(\hat{P}_1\) for examining \(t_0\) is \(O(K^3)\) [27], the overall complexity of the method can be calculated as \(O(\log(\frac{1}{\epsilon})K^2)\). In order to further analyze single-beam beamspace MIMO-NOMA in terms of its max-min rate, we will carry out simulations in the next subsection, and show that its max-min rate always grows rapidly with respect to SNR.

C. Simulations and Analysis

In this subsection, simulations are carried out for the max-min rate of the beamspace MIMO-NOMA and the beamspace MIMO-OMA in the single-beam case. For the fairness of comparison, the same beam selection scheme is adopted in the two systems. Particularly, beamspace MIMO-OMA adopts the classic OMA scheme of time division multiple access (TDMA) to serve multiple users. Since only one RF chain is deployed at the BS, only one user can be served in each time slot, and the number of time slots is equal to the number of the users. To maximize the minimal rate, the user will be served with full power in each time slot [22]. For both of the single-beam systems, \(N = 64\) antennas and \(N_{\text{RF}} = 1\) RF chain are deployed at the BS, while 1 LoS component and \(L = 5\) NLoS components are assumed for the mmWave channel. The complex gain \(\beta_k^{(0)}\) of the LoS component for an arbitrary \(k\)-th user follows complex Gaussian distribution \(CN(0,1)\), and that of the NLoS component \(\beta_k^{(l)}\) follows \(CN(0,1)\), with \(l\) varying from 1 to \(L\) [23]. The spatial directions \(\psi_k^{(0)}\) and \(\psi_k^{(l)}\) for the LoS and NLoS components follow uniform distribution within \([-\frac{\pi}{2}, \frac{\pi}{2}]\), respectively [23]. Additionally, the simulation results are obtained using CVX [32].

Fig. 2 presents the max-min rates of the two single-beam systems serving 16 users with respect to the SNR, which is defined as \(\log_{10}(P_{\text{max}}/\sigma^2)\). It shows that the max-min rates of both systems increase rapidly with the SNR, and the growing slopes of the two corresponding curves almost remain unchanged. This is because single-beam beamspace MIMO-OMA is interference-free, and also because intra-beam interference of the single beam beamspace MIMO-NOMA has no restriction on its max-min rate, which will be detailedly analyzed in Section V. As a result, in the single-beam case, beamspace MIMO-NOMA is superior to beamspace MIMO-OMA in terms of the max-min rate in the whole SNR region, which is because NOMA can efficiently utilize the limited communication resources.

Besides, Fig. 3 depicts the relationship between the two max-min rates and the number of served users, with the SNR set to 18 dB. From this figure, we can see that when the number of served users increases, both of the max-min rates decrease due to the limited maximum transmit power. In addition, since NOMA is adopted, the max-min rate of the beamspace MIMO-NOMA is always larger than that of the beamspace MIMO-OMA in the single-beam case.

IV. MAX-MIN FAIRNESS ANALYSIS FOR MULTI-BEAM BEAMSPACE MIMO-NOMA

In this section, we will conduct the analysis for the multi-beam case, where multiple RF chains are deployed at the BS. Particularly, the difference between the single-beam case and the multi-beam case in terms of the inter-beam interference is pointed out, and a more complicated multi-beam power allocation problem is formulated to maximize the minimal rate, which is also solved by a bisection approach. Moreover, our analysis reveal that beamspace MIMO-NOMA is not always superior to the beamspace MIMO-OMA in the multi-beam case, especially when the inter-beam interference is relatively
severe, which is owing to the high SNR or the large number of served users.

A. Problem Formulation

The minimal rate of beamspace MIMO-NOMA in the multi-beam case is maximized by formulating a multi-beam power allocation problem. Note that within each beam, the SIC decoding order is based on the order of effective channel qualities, i.e., \(|\|\hat{h}_{m,1}^H w_m\|^2 > |\|\hat{h}_{m,2}^H w_m\|^2 > \cdots > |\|\hat{h}_{m,[S_m]}^H w_m\|^2|\). According to equations (23), we cannot decide whether \(m\)-th user’s signal, and the corresponding SINR \(\gamma_{n,i}^m\) can be expressed as:

\[
\gamma_{n,i}^m = \frac{\|\hat{h}_{m,i}^H w_m\|^2}{\varepsilon_{n,i}^m + \sigma^2 p_{m,n}}.
\]

In the above equation (22), \(\varepsilon_{n,i}^m\) denotes the power of the intra and inter-beam interferences:

\[
\varepsilon_{n,i}^m = \sum_{n' = 1}^{n-1} \|\hat{h}_{m,i}^H w_m\|^2 p_{m,n'} + \sum_{m' \neq m, n' = 1}^{[S_m]} \|\hat{h}_{m',i}^H w_m\|^2 p_{m',n'}. \tag{23}
\]

It is worth pointing out that due to the inter-beam interference term presented in equation (23), we cannot decide whether \(\gamma_{n,i}^m\) is larger than \(\gamma_{n,j}^m\) \((i < j)\) or not before power allocation:

\[
\gamma_{n,i}^m - \gamma_{n,j}^m = \frac{p_{m,n} \left(\|\hat{h}_{m,i}^H w_m\|^2 - \|\hat{h}_{m,j}^H w_m\|^2\right)}{\left(\varepsilon_{n,i}^m + \sigma^2\right)\left(\varepsilon_{n,j}^m + \sigma^2\right)}.
\]

in equation (24) is manipulated from the inter-beam interference term in equation (23), and its sign cannot be determined without the knowledge of power allocation. Therefore, unlike the single-beam case, it is unknown which user decides the rate for the \(n\)-th user in the \(m\)-th beam before power allocation, which is mainly due to the inter-beam interference. This is a significant difference between the single-beam case and the multi-beam case. Based on the above discussions, the objective function to maximize the minimal rate can be written as

\[
\max \min_{\{p_{m,n}\}} \{R_{n,i}^m\} = \max \min_{\{p_{m,n}\}} \{R_{n,i}^m\}. \tag{25}
\]

Besides, the two power constraints are similar to (15) and (16) in Section III-A. Therefore, the formulated multi-beam power allocation problem is given by

\[
P_2 : \max \min_{\{p_{m,n}\}} \{R_{n,i}^m\}, \tag{26}
\]

s.t. \(C_1 : p_{m,n} \geq 0, \quad \forall m, n,\)

\[
C_2 : \sum_{m=1}^{NRF} \sum_{n=1}^{[S_m]} ||w_m||^2 \leq P_{\text{max}},
\]

and the maximized minimal rate is denoted as \(r^*\).

Similar to its single-beam counterpart, the multi-beam power allocation problem \(P_2\) is also non-convex. Moreover, due to the inter-beam interference, the achievable rate for each user has to be expressed by the minimum of a series of rates, which makes the multi-beam problem more complex. In order to solve this non-convex optimization problem, a bisection approach will be introduced in detail in Section IV-B.

B. Minimal Rate Maximization

Since we aim at maximizing the minimal rate of the multi-beam beamspace MIMO-NOMA system, the existing algorithms used to maximize the sum rate (such as WMMSE) cannot be applied to solve our problem. To this end, a bisection approach is utilized in this subsection to maximize the minimal
rate for the system. To deal with the complicated objective function in (26), an auxiliary variable \( t \) is introduced to equivalently transform the original objective function into constraint \( C_3 \) of the following new problem \( P'_2 \):

\[
P'_2 : \max_{(p_{m,n})} t,
\]

\[
s.t. \ C_1: p_{m,n} \geq 0, \ \forall m,n,
\]

\[
C_2: \sum_{m=1}^{N_{RF}} |S_m| \sum_{n=1}^{N_{SF}} \|\mathbf{w}_m\|^2 p_{m,n} \leq P_{\text{max}},
\]

\[
C_3: R_{m,i}^n \geq t, \ \forall m,n, \ 1 \leq i \leq n.
\]

(27)

Inspired by Algorithm 1 in Section III-B, we set the auxiliary variable \( t \) to some \( t_0 \) in a bisection manner, and construct the following power consumption minimization problem to examine whether \( t_0 \) is larger than \( r^* \) or not:

\[
\hat{P}_2 : \min_{(p_{m,n})} \sum_{m=1}^{N_{RF}} \sum_{n=1}^{N_{SF}} \|\mathbf{w}_m\|^2 p_{m,n},
\]

\[
s.t. \ C_1: p_{m,n} \geq 0, \ \forall m,n,
\]

\[
C_2: R_{m,i}^n \geq t_0, \ \forall m,n, \ 1 \leq i \leq n.
\]

(28)

As mentioned before, due to the inter-beam interference, it is unknown which user decides the rate for the \( n \)-th user in the \( m \)-th beam before power allocation. Therefore, unlike the single-beam case, equations cannot be derived from problem \( \hat{P}_2 \), and thus closed-form expressions cannot be directly deduced. Fortunately, the objective function as well as the constraint \( C_1 \) of the problem \( \hat{P}_2 \) are linear with respect to the power allocation parameters \( \{p_{m,n}\} \), while the constraint \( C_2 \) can be shown to be linear by the following equation (29):

\[
R_{m,i}^n \geq t_0
\]

\[
\iff \|\mathbf{h}_{m,i}^H \mathbf{w}_m\|^2 \geq \sum_{k=1}^{n-1} \|\mathbf{h}_{m,k}^H \mathbf{w}_j\|^2 \sum_{j \neq m} |S_j| + \left( \sum_{k=1}^{n-1} \sum_{j \neq m} \|\mathbf{h}_{m,k}^H \mathbf{w}_j\|^2 \right) p_{j,k} + \sum_{j \neq m} \|\mathbf{h}_{m,j}^H \mathbf{w}_j\|^2 p_{j,m,n} - \|\mathbf{h}_{m,i}^H \mathbf{w}_m\|^2 \geq 2t_0 - 1 + \sigma^2 \leq 0.
\]

(29)

Therefore, the problem \( \hat{P}_2 \) can also be solved by classic linear programming methods. Denote the minimized power consumption for problem \( \hat{P}_2 \) as \( P^* \), we also have \( r^* \geq t_0 \) when \( P^* \leq P_{\text{max}} \), and \( r^* < t_0 \) when \( P^* > P_{\text{max}} \). Correspondingly, the bisection-based method is presented in Algorithm 2. Similar to its single-beam counterpart, the bisection-based method in the multi-beam case can also converge to the optimal solution with a desirable accuracy \( \epsilon \) with linear convergence [31], which demonstrates its effectiveness, and the number of required iterations is also \( O(\log(\frac{t_0}{\epsilon})) \). Since in each iteration the computation complexity to solve \( \hat{P}_2 \) for examining \( t_0 \) is \( O(K^{3.5}) \) using classic linear programming methods [33], its complexity can be calculated as \( O(\log(\frac{t_0}{\epsilon})K^{3.5}) \).

Based on the bisection procedure presented above, in the following paragraphs, we will prove that max-min rate is achieved when the rates of NOMA users are equal to one another for multi-beam systems. Particularly, for an arbitrary \( n \)-th user in the \( m \)-th beam, its signal must be decoded by the \( i \)-th user for all \( i \leq n \), and thus its rate is decided by a certain \( j \)-th user for \( j \leq n \). This \( j \)-th user is unknown to us as mentioned before, but it could be the first user, the second user, etc. Accordingly, the power consumption minimization problem \( \hat{P}_2 \) can be decomposed into a set of problems \( \hat{P}_\lambda \) by traversing all the possibilities. Since each beam can accommodate \( |S_m| \) users, the total number of problems in set \( \hat{P}_\lambda \) can be calculated by \( \prod_{m=1}^{N_{RF}} |S_m|! \). To be more specific, a certain power consumption minimization problem \( \hat{P}_3 \) in set \( \hat{P}_\lambda \) can be formulated as

\[
\hat{P}_3 : \min_{(p_{m,n})} \sum_{m=1}^{N_{RF}} \sum_{n=1}^{N_{SF}} \|\mathbf{w}_m\|^2 p_{m,n},
\]

\[
s.t. \ C_1: p_{m,n} \geq 0, \ \forall m,n,
\]

\[
C_2: R_{m,j}^{m,n} \geq t_0, \ \forall m,n.
\]

(30)

where \( R_{m,j}^{m,n} \) indicates that the rate for the \( n \)-th user in the \( m \)-th beam is determined by the \( j(m,n) \)-th user in the same beam. We can denote the set of solutions for all the problems in set \( \hat{P}_\lambda \) as \( \hat{S}_\lambda \), and the solution for problem \( \hat{P}_2 \) must lie in set \( \hat{S}_\lambda \). Then, we prove that for every problem in set \( \hat{P}_\lambda \), its solution satisfies the property that the rates of NOMA users are equal to one another.

**Lemma 1:** For an arbitrary problem \( \hat{P}_3 \) in set \( \hat{P}_\lambda \), its solution satisfies the property that the rates of all NOMA users are equal to one another.

**Proof:** Since problem \( \hat{P}_3 \) is decomposed from \( \hat{P}_2 \), it is also convex, and KKT condition can be utilized to prove this lemma. Particularly, for an arbitrary \( n \)-th user in the \( m \)-th beam, the corresponding power allocation parameters satisfy the KKT conditions presented below:

\[
(t_0^2 - 1) \left( \sum_{m'=n+1}^{N_{RF}} \lambda_{m,m'} \left\| \mathbf{h}_{m,j(m,m')}^H \mathbf{w}_m \right\|^2_2 \right)
\]

\[
+ \sum_{m' \neq n} \lambda_{m',m} \left( \left\| \mathbf{h}_{m',j(m,m')}^H \mathbf{w}_{m'} \right\|^2_2 + \left\| \mathbf{w}_m \right\|^2_2 \right) = \lambda_{m,n} \left( \left\| \mathbf{h}_{m,j(m,m)}^H \mathbf{w}_m \right\|^2_2 + \mu_{m,n} \right)
\]

(31)
\begin{align}
(2^{t_0} - 1)(\xi_{m,j(m,n)}^m + \sigma^2) & \leq \|H_{m,j(m,n)}^m w_m\|_2^2 p_{m,n} \tag{32} \\
\lambda_{m,n} \left( (2^{t_0} - 1)(\xi_{m,j(m,n)}^m + \sigma^2) \right. & \left. - \|H_{m,j(m,n)}^m w_m\|_2^2 p_{m,n} \right) = 0 \tag{33} \\
p_{m,n} & \geq 0, \lambda_{m,n} \geq 0, \mu_{m,n} \geq 0, \tag{34}
\end{align}

where
\[
\xi_{m,j(m,n)}^m = \sum_{n' = 1}^{n-1} \|H_{m,j(m,n)}^m w_m\|_2^2 p_{m,n'} + \sum_{n' \neq m} \sum_{n'' = 1}^{n} \|H_{m,j(m,n)}^m w_{m'}\|_2^2 p_{m',n'}
\]

denotes the power for intra-beam and inter-beam interferences. Since the left-hand-side of equation (32) is positive, the corresponding right-hand side is also positive, which indicates that \(p_{m,n} > 0\). Thus, according to equation (34), we have \(\mu_{m,n} = 0\). Moreover, since the left-hand-side of equation (IV-B) is positive, its right-hand side must also be positive. Given that \(\mu_{m,n} = 0\), we now have \(\lambda_{m,n} > 0\).

Based on the above discussions, and note that equation (IV-B) is the complementary slackness condition for constraint \(C_2\) of the problem \(P_3\), we can now conclude that the solution must satisfy constraint \(C_2\) with equality. In this way, the rates for NOMA users will all be equal to \(t_0\), and this completes the proof.

With \textit{lemma 1} proved above, now we can prove that the optimal solution for the original multi-beam power allocation problem \(P_2\) satisfies the property that the rates of all the NOMA users are equal to one another.

\textbf{Theorem 1:} The max-min rate of the original multi-beam power allocation problem \(P_2\) can be achieved when all the rates are assigned the same rate.

\textbf{Proof:} Based on \textit{lemma 1}, we can conclude that each solution in set \(S_C\) satisfies the property that the rates of all NOMA users are equal to one another. Since the solution for problem \(P_2\) lies in set \(S_C\), the solution for \(P_2\) will also satisfy this property. Finally, given that the max-min rate of the original multi-beam power allocation problem \(P_2\) is obtained by solving a series of problems \(P_2\), it will be achieved when the rates for all NOMA users are equal to one another [27], which completes the proof.

Theorem 1 shows that when we optimize the sum rate of the system with the minimum transmission rate equal to the max-min rate, then the rate for all the NOMA users must be equal to the max-min rate. At this time, there will be no loss in terms of the sum rate between optimizing the minimal rate and optimizing the sum rate. Moreover, when the minimum transmission rate decreases from the max-min rate, the power allocated to the weak user will become smaller, and the optimized sum rate of the system will gradually become larger. Therefore, the sum rate loss between optimizing the minimal rate and optimizing the sum rate will also become larger.

In order to further analyze the performance of the multi-beam beamspace MIMO-NOMA in terms of the max-min rate, we will carry out simulations in the next subsection, and show that the growing slope of its max-min rate obviously decreases in the high-SNR region.

\section{Simulations and Analysis}

In this subsection, simulations are carried out for the max-min rate of the beamspace MIMO-NOMA and the beamspace MIMO-OMA in the multi-beam case. For the fairness of comparison, the same beam selection scheme is adopted in the two systems. Similar to that in Section III-C, multiple users in each beam of the beamspace MIMO-OMA are served by TDMA, and the number of time slots is equal to the maximum number of users within each beam [24]. Since the BS is equipped with multiple RF chains, multiple users will be served in each time slot. To maximize the minimal rate, for each time slot, zero-forcing is adopted to eliminate interference, and power allocation is optimized using the similar bisection procedure in this paper. Besides, the BS is equipped with \(N = 64\) antennas, and the number of RF chains \(N_{RF}\) varies from 4 to 8, while the other simulation settings are the same as that in Section III-C.

Fig. 4 presents the max-min rates of the two multi-beam systems serving 16 users with respect to the SNR, where two scenarios are considered with the number of RF chains set to 4 and 8, respectively. It can be shown from the figure that the max-min rate with 8 RF chains is larger than the max-min rate with 4 RF chains. In fact, with the increase of the number of RF chains, the number of beams will also increase. Therefore, the number of users served by each beam will decrease, which mitigates the intra-beam interference. As a result, the max-min rate with more RF chains will be larger than that with less RF chains. Besides, this figure also shows that for both of the scenarios, with the increase of the SNR, the growing slope for the max-min rate of the beamspace MIMO-NOMA obviously decreases in the high-SNR region, while that of the beamspace MIMO-OMA almost remains unchanged. As a result, the max-min rate of the beamspace MIMO-NOMA becomes smaller than that of the beamspace MIMO-OMA in the high-SNR region, while the former is actually larger than the latter in the low-SNR region, which is shown by the intersection points of the two corresponding curves.

To explain this phenomenon, we would like to point out that beamspace MIMO-NOMA is rank-deficient, since the number of the RF chains \(N_{RF}\) is smaller than that of the served users \(K\). Thus, the inter-beam interference of beamspace MIMO-NOMA in the multi-beam case cannot be completely eliminated by digital precoding, which largely constrains the growing slope of its max-min rate and makes it interference-limited in the high-SNR region. In contrast, beamspace MIMO-OMA is free of interference due to the TDMA scheme and the zero-forcing method. Therefore, the max-min rate of multi-beam beamspace MIMO-OMA will finally exceed that of the multi-beam beamspace MIMO-NOMA with sufficiently large SNR, which indicates that the intersection point always exists. Note that the intersection
point of the two multi-beam systems with 4 RF chains appears at a larger SNR than that with 8 RF chains, and this is because the inter-beam interference of the former is less severe than the latter due to smaller number of beams.

With the SNR set to 18 dB, Fig. 5 depicts the max-min rates of the two systems against the number of served users with \( N_{RF} \) set to 4 and 8, respectively. It can be observed from the figure that since the inter-beam interference gets increasingly severe when the number of served users increases, beamspace MIMO-NOMA almost has no advantage over beamspace MIMO-OMA in terms of the max-min rate when serving a large number of users.

So far, we have conducted max-min fairness analysis for beamspace MIMO-NOMA in both the single-beam case and the multi-beam case. Specifically, although affected by the intra-beam interference, the max-min rate of the single-beam beamspace MIMO-NOMA is not interference-limited, and is always larger than that of the single-beam beamspace MIMO-OMA, as shown in Fig. 2. However, the max-min rate of the multi-beam beamspace MIMO-NOMA affected by the inter-beam interference is interference-limited, and is smaller than that of the multi-beam beamspace MIMO-OMA in high-SNR region, as shown in Fig. 4. To explain this difference, we will point out that the interfering patterns of intra-beam and inter-beam interferences are completely different in the next section.

V. DIFFERENT PATTERNS OF INTRA-BEAM AND INTER-BEAM INTERFERENCES

In this section, we will comprehensively analyze the different interfering patterns of the intra-beam and inter-beam interferences, and point out their different effects on the max-min rate. Particularly, the intra-beam interference is of successive pattern, which is proved to have no restriction on the max-min rate. However, the inter-beam interference is of mutual pattern, which is proved to restrict the max-min rate to a derived upper bound.

A. Successive Pattern of the Intra-Beam Interference

In the single-beam beamspace MIMO-NOMA system, the \( n \)-th user's achievable rate \( R_n \) is actually determined by the rate of decoding its signal by itself, as shown by equations (12) and (13):

\[
R_n = \log_2 \left( 1 + \frac{|\hat{h}_{nn}|^2 p_n}{\sum_{k=1}^{n-1} |\hat{h}_{nk}|^2 p_k + \sigma^2} \right). \quad (36)
\]

As shown by the denominator of the above equation (36), the successive pattern of the intra-beam interference means that the interference of the \( n \)-th user only comes from the users with better effective channel qualities, i.e., users whose indices are smaller than \( n \). This is because SIC will successively decode and remove the signals of the users with worse effective channel qualities [27]. Therefore, assume that the \( n \)-th user is interfered by the \( k \)-th user for \( k < n \), then the \( k \)-th user is not interfered by the \( n \)-th user, since the interference of the \( k \)-th user only comes from the users with indices smaller than \( k \). This indicates that the users will not mutually interfere with one another in the single-beam case. As a result, we will show that the intra-beam interference will not limit the max-min rate by proving the following Proposition 1:

**Proposition 1:** In the single-beam beamspace MIMO-NOMA system, the max-min rate \( r^* \) will not be limited by the intra-beam interference, and can achieve any arbitrarily large value \( r_0 \) as long as the total transmit power \( P \) is sufficient.

**Proof:** We will prove this proposition by showing that the achievable rate of each NOMA user in the single-beam beamspace MIMO-NOMA system can achieve any arbitrarily large \( r_0 \) with sufficiently large \( P \). Particularly, we plan to firstly set the achievable rates for all NOMA users equal to the arbitrarily large \( r_0 \), and examine whether this can be achieved by calculating the corresponding power allocation parameters \( \{p_n\} \). Thus, the following set of equations are formulated:

\[
R_n = r_0 \Leftrightarrow \sum_{k=1}^{n-1} p_k - \frac{p_n}{2r_0 - 1} = 0. \quad (37)
\]
From the above equation (37), we have
\[ p_n = (2^{r_o} - 1) \left( \sum_{k=1}^{n-1} p_k + \frac{\sigma^2}{|h_{m_1}^r|^2} \right). \tag{38} \]

Therefore, once the power allocation parameters \( \{p_k\}_{k=1}^{n-1} \) are obtained, the power allocation parameter \( p_n \) can be calculated by the above equation (38). Note that \( p_1 \) can be directly calculated by \( p_1 = (2^{r_0} - 1) \sum_{k=1}^{n} \frac{1}{|h_{m_1}^r|^2} \), and thus we can obtain the other power allocation parameters \( p_2, p_3, \ldots, p_K \) successively with (38). Moreover, it is straightforward to find that all the power allocation parameters are no less than zero. Since the corresponding power allocation parameters can be calculated, the achievable rate of each NOMA user can achieve any arbitrarily large \( r_0 \), and thus the max-min rate \( r^* \) of all NOMA users can achieve \( r_0 \), which completes the proof.

\[ \blacksquare \]

B. Mutual Pattern of the Inter-Beam Interference

In the multi-beam case, the inter-beam interference is of mutual pattern, which refers to the fact that different users served by different beams interfere with one another. To be more specific, take the \( n_1 \)-th user in the \( m_1 \)-th beam as an example, according to (8), the power of all the interferences for decoding its signal at itself is given by
\[ c_{n_1, n_1}^{m_1} = \sum_{k=1}^{n_1-1} \left| h_{m_1, n_1}^{H} w_{m_1} \right|^2 p_{m_1,k} + \left| \sum_{k=1}^{n_2} \left| h_{m_1, n_1}^{H} w_{m_2} \right|^2 p_{m_2,k} \right| + \sum_{j \neq m_1, m_2} \left| \sum_{k=1}^{n} \left| h_{m_1, n_1}^{H} w_j \right|^2 p_{j,k} \right|, \tag{39} \]
where \( m_2 \neq m_1 \) is the index for another beam. Besides, the power of interference caused by users in the \( m_2 \)-th beam is denoted by the second term of the right hand side in (39). This indicates that an arbitrary \( n_2 \)-th user in the \( m_2 \)-th beam will interfere with the \( n_1 \)-th user in the \( m_1 \)-th beam, and vice versa, which accounts for the mutual pattern of the inter-beam interference. Therefore, even if we only focus on these two arbitrary users mentioned above, the following Proposition 2 shows that the inter-beam interference between them will largely constrain the max-min rate, and the corresponding upper bound is also derived:

**Proposition 2**: In the multi-beam beamspace MIMO-NOMA system, for any \( m_1, m_2, n_1, \) and \( n_2 \) satisfying
\[ m_1 \neq m_2, 1 \leq m_1, m_2 \leq N_{RF}, \]
\[ 1 \leq n_1 \leq |S_{m_1}|, 1 \leq n_2 \leq |S_{m_2}|, \tag{40} \]
we have
\[ r^* < \max \left\{ B_{m_1, n_1}^{m_2, n_2}, B_{m_2, n_2}^{m_1, n_1} \right\}, \tag{41} \]
where
\[ B_{m_1, n_1}^{m_2, n_2} = \log_2 \left( 1 + \frac{\left| h_{m_2, n_2}^{H} w_{m_2} \right|^2}{\left| h_{m_1, n_1}^{H} w_{m_1} \right|^2} \right), \tag{42} \]
and
\[ B_{m_2, n_2}^{m_1, n_1} = \log_2 \left( 1 + \frac{\left| h_{m_1, n_1}^{H} w_{m_1} \right|^2}{\left| h_{m_2, n_2}^{H} w_{m_2} \right|^2} \right). \tag{43} \]

**Proof**: We will prove this proposition by showing that for any power allocation parameters \( \{p_{m,n}\} \), the minimal rate of all NOMA users \( r_{min} \) satisfies \( r_{min} < \max \left\{ B_{m_1, n_1}^{m_2, n_2}, B_{m_2, n_2}^{m_1, n_1} \right\} \).

Firstly, since the overall minimum must be no larger than the local minimum, the minimal rate of all NOMA users must be no larger than the minimal rate of two arbitrary users. Therefore, according to (25), we have
\[ r_{min} = \min_{m,n} \left\{ R_{n}^{m} \right\} = \min_{m,n, 1 \leq i \leq n} \left\{ R_{n_i}^{m} \right\} \leq \min_{m,n} \left\{ R_{n}^{m} \right\} \leq \min \left\{ R_{m_1, n_1}^{m_2, n_2}, R_{m_2, n_2}^{m_1, n_1} \right\}. \tag{44} \]

To elaborate a little further, \( R_{n_1}^{m_1} \) is the achievable rate for user \( n_1 \) in beam \( m_1 \) to decode its own signal:
\[ R_{n_1}^{m_1} = \log_2 \left( 1 + \frac{\left| h_{m_1, n_1}^{H} w_{m_1} \right|^2}{\left| h_{m_2, n_2}^{H} w_{m_2} \right|^2} \left( p_{m_1, n_1} \right) \right). \tag{45} \]

For this user, its interference includes intra-beam interference within the same beam and inter-beam interference from all the users in every other beam. Therefore, the interference power must be larger than the power of the interference only coming from user \( n_2 \) in beam \( m_2 \):
\[ \xi_{n_1, n_1}^{m_1} = \sum_{k=1}^{n_1-1} \left| h_{m_1, n_1}^{H} w_{m_1} \right|^2 p_{m_1,k} + \sum_{k=1}^{n_2} \left| h_{m_1, n_1}^{H} w_{m_2} \right|^2 p_{m_2,k} + \sum_{j \neq m_1, m_2} \sum_{k=1}^{n} \left| h_{m_1, n_1}^{H} w_j \right|^2 p_{j,k} > \left| h_{m_1, n_1}^{H} w_{m_2} \right|^2 p_{m_2,n_2}. \tag{46} \]

Combining (45) and (46), it holds that
\[ R_{n_1}^{m_1} = \log_2 \left( 1 + \frac{\left| h_{m_1, n_1}^{H} w_{m_1} \right|^2}{\left| h_{m_1, n_1}^{H} w_{m_2} \right|^2} \left( p_{m_1, n_1} \right) \right) < \log_2 \left( 1 + \frac{\left| h_{m_1, n_1}^{H} w_{m_1} \right|^2}{\left| h_{m_1, n_1}^{H} w_{m_2} \right|^2} \left( p_{m_2, n_2} \right) \right). \tag{47} \]

Similarly, for user \( n_2 \) in beam \( m_2 \), it also holds that
\[ R_{n_2}^{m_2} < \log_2 \left( 1 + \frac{\left| h_{m_2, n_2}^{H} w_{m_2} \right|^2}{\left| h_{m_2, n_2}^{H} w_{m_1} \right|^2} \left( p_{m_2, n_2} \right) \right) \tag{48} \]

\[ \blacksquare \]
Based on (47), when \( p_{m_1,n_1} \leq p_{m_2,n_2} \), we have,

\[
\min \{ R_{m_1,n_1}^{m_1}, R_{m_2,n_2}^{m_2} \} \\
\leq R_{m_1,n_1}^{m_1} \\
< \log_2 \left( 1 + \frac{\| H_{m_1,n_1}^H w_{m_1} \|_2^2}{\| H_{m_1,n_2}^H w_{m_2} \|_2^2} \right) \\
= B_{m_1,n_1}^{m_1},
\]

and based on (48) it holds that,

\[
\min \{ R_{m_1,n_1}^{m_1}, R_{m_2,n_2}^{m_2} \} \\
\leq R_{m_2,n_2}^{m_2} \\
< \log_2 \left( 1 + \frac{\| H_{m_2,n_2}^H w_{m_2} \|_2^2}{\| H_{m_2,n_2}^H w_{m_1} \|_2^2} \right) \\
= B_{m_2,n_2}^{m_2},
\]

(49)

when \( p_{m_1,n_1} > p_{m_2,n_2}. \) With the help of the equations (49) and (50) mentioned above, we can now claim that no matter \( p_{m_1,n_1} \leq p_{m_2,n_2} \) or \( p_{m_1,n_1} > p_{m_2,n_2}, \) it always holds that

\[
\min \{ R_{m_1,n_1}^{m_1}, R_{m_2,n_2}^{m_2} \} < \max \{ B_{m_2,n_2}^{m_2}, B_{m_1,n_1}^{m_1} \}. \quad (51)
\]

According to (44) and (51), we have

\[
r_{\min} \leq \min \{ R_{m_1,n_1}^{m_1}, R_{m_2,n_2}^{m_2} \} \\
< \max \{ B_{m_2,n_2}^{m_2}, B_{m_1,n_1}^{m_1} \},
\]

(52)

for any power allocation parameters \( \{ p_{m,n} \} \), and thus

\[
r^* = \max \{ p_{m,n} \} r_{\min} < \max \{ B_{m_2,n_2}^{m_2}, B_{m_1,n_1}^{m_1} \}. \quad (53)
\]

This completes the proof.

It is worth pointing out that due to the mutual pattern of the inter-beam interference between user \( n_1 \) in beam \( m_1 \) and user \( n_2 \) in beam \( m_2 \), the two corresponding rates \( R_{m_1,n_1}^{m_1} \) and \( R_{m_2,n_2}^{m_2} \) are upper-bounded by \( B_{m_1,n_1}^{m_1} \) and \( B_{m_2,n_2}^{m_2} \), respectively, as shown by (49) and (50), which finally leads to the upper bound for the max-min rate \( r^* \). Moreover, note that the two terms \( \| H_{m_1,n_1}^H w_{m_1} \|_2^2 \) and \( \| H_{m_1,n_2}^H w_{m_2} \|_2^2 \) are on the denominator of \( B_{m_2,n_2}^{m_2} \) and \( B_{m_1,n_1}^{m_1} \), respectively, and they denote the interference from the \( m_1 \)-th beam to the \( m_2 \)-th beam and vice versa. They show that the max-min rate is mainly restricted by the inter-beam interference. Specifically, when the inter-beam interference is more severe, the max-min rate will become more restricted, and the upper bound will become smaller. Besides, since the multi-beam beamspace MIMO-OMA is free of interference, its max-min rate will finally exceed that of the multi-beam beamspace MIMO-NOMA with sufficiently large SNR, which accounts for the intersection point in Fig. 4.

In addition, the upper-bound obtained in Proposition 2 is only relevant to a pair of users served by two different beams. Thus, once two arbitrary users in different beams are selected, we can utilize Proposition 2 to calculate one corresponding upper bound, and a set \( \mathcal{U} \) of the upper bounds can be constructed in this way:

\[
\mathcal{U} = \{ \max \{ B_{m_1,n_1}^{m_1}, B_{m_2,n_2}^{m_2} \} | m_1, m_2, n_1, n_2 \text{ satisfiying } (40) \}.
\]

With the set \( \mathcal{U} \) constructed as mentioned above, a more tighter upper bound \( r_{\text{bound}} \) can be obtained by searching for the minimum in set \( \mathcal{U} \):

\[
r_{\text{bound}} = \min \{ \max \{ B_{m_1,n_1}^{m_1}, B_{m_2,n_2}^{m_2} \} \},
\]

s.t. \( \max \{ B_{m_1,n_1}^{m_1}, B_{m_2,n_2}^{m_2} \} \in \mathcal{U} \). \quad (54)

Using the formulas of combination, the number of elements in set \( \mathcal{U} \) can be readily calculated as:

\[
|\mathcal{U}| = C_K^2 - \sum_{m=1}^{N_{\text{RF}}} C_{S_m}^2.
\]

Therefore, the minimum in set \( \mathcal{U} \) can always be obtained by searching among \( |\mathcal{U}| \) elements. In the next subsection, we will verify the derived upper bound \( r_{\text{bound}} \) in (54) via simulations.

C. Simulations and Analysis

We conduct simulations in this subsection to verify the derived upper bound for the max-min rate of the multi-beam beamspace MIMO-NOMA, while maintaining the same simulation settings as that in Section IV-C. Particularly, two scenarios with \( N_{\text{RF}} \) varying from 4 to 8 are considered. Fig. 6 presents the derived upper bound calculated by (54) and the
max-min rate with respect to the SNR for $N_{\text{RF}} = 4$, while those for $N_{\text{RF}} = 8$ are shown in Fig. 7. Both of the two figures show that the max-min rate nearly stops growing in the high-SNR region, which demonstrates the effect of the inter-beam interference. Moreover, the max-min rate is always smaller than the derived upper bound, which verifies our theoretical analysis in Section V-B. Note that the upper bound is larger for smaller number of RF chains. This is because with smaller number of RF chains, the number of beams is also smaller, and the inter-beam interference will become less severe. As mentioned in Section V-B, when the inter-beam interference is less severe, the max-min rate will become less restricted, and the upper bound will become larger.

VI. CONCLUSION

In this paper, we have maximized the minimal rate of beamspace MIMO-NOMA in both the single-beam and multi-beam cases. Particularly, in the single-beam case, we find that the maximized minimal rate of beamspace MIMO-NOMA grows rapidly in the whole SNR region, while in the multi-beam case grows slower and slower and becomes interference-limited in the high-SNR region. To explain this difference, we have comprehensively analyzed the different interfering patterns of the intra-beam and inter-beam interfences, pointed out their different effects on the maximized minimal rate, and derived the upper bound for the multi-beam case. More importantly, when considering max-min rate in the multi-beam case, we have demonstrated that beamspace MIMO-NOMA is superior to beamspace MIMO-OMA in the low SNR region, while the former becomes inferior to the latter in the high-SNR region. Note that the above analysis does not deny the application of NOMA, but points out that NOMA also has its weaknesses in spite of its known strengths in terms of the connection number, sum capacity, etc., which completes the analysis framework and provides a deeper understanding for NOMA.

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